Faculty of Electrical Engineering, Mathematics, and Computer Science Delft University of Technology

exam – **Embedded Software** – TI2726-B January 29, 2018 13.30 - 15.00

This exam (6 pages) consists of 60 True/False questions. Your score will be computed as: $max(0, \frac{\#correct}{60} - \frac{1}{2}) \times 2 \times 9 + 1$ It is **not** allowed to consult the book, handouts, or any other notes.

Instructions for filling in the answer sheet:

- You may use a **pencil** (erasures are allowed) or a **pen** (blue or black, **no** red, **no** strike outs).
- Fill in the boxes **completely**.
- Answer **all** questions; there is no penalty for guessing.
- Do not forget to fill in your Name and Student Number, and to sign the form.

The following abbreviations are assumed to be known:

- RR (Round Robin)
- RRI (Round Robin with Interrupts)
- FQS (Function Queue Scheduling)
- RTOS (Real-Time Operating System)
- ISR (Interrupt Service Routine)
- UART (Universal Asynchronous Receiver Transmitter)

One system clock tick = 10 ms (unless stated otherwise).

We make use of the following definitions:

```
void delay(int ms) {
    !! do some CPU computation to the number of ms milliseconds
}
void putchar(char c) {
    while (!! UART tx buffer not empty)
    ;
    !! send c to UART tx buffer
}
void puts(char *s) {
    !! write string s using putchar
}
```

1.	Embedded programming is more difficult than "classical" programming because of the lack of support for recursion.	true/false
2.	A defining characteristic of embedded systems is the restricted, or complete lack, of a user interface.	true/false
3.	Several models of computation for embedded systems are described in [Lee:2002]. - The ROS software (used in the practicals) is a prime example of the publish-and- subscribe model.	true/false
4.	The Underground Tank Monitoring System is a somewhat contrived example of an embedded system as it involves input (sensors/buttons) and output (display/printer), but lacks real-time constraints and resource limitations.	true/false
5.	Despite advances in software engineering practices, as a rule of thumb, embedded software contains 1-10 bugs per thousand lines of code.	true/false
6.	Hardware interrupts can be disabled; software interrupts cannot.	true/false
7.	An embedded program can be coded as a finite state machine; the number of incoming transitions (arcs) into a state S must equal the number of outgoing transitions (arcs).	true/false
8.	Finite State Machines can be coded in a number of ways in C.In the table-based solution, every transition (arc) is encoded as a separate function.	true/false
9.	<pre>int main(void) { int c; statefp state = start; while((c = getchar()) != EOF) { state = (statefp) (*state)(c); } return 0; }</pre>	
	The above loop drives the FSM until the end state is reached.	true/false
10.	Unlike recursive data structures, recursive function types cannot be properly defined in C and require kludges like void pointers and type casts.	true/false
11.	The C language does not contain a built-in type to represent booleans. - in control flow statements, expressions evaluating to 0 are regarded as logically False.	true/false
12.	<pre>typedef void (* resolve)(void *old, void *new);</pre>	
	The first pair of parenthesis in the definition above is for clarity (stressing a function pointer is involved) and can be left out without changing the meaning.	true/false
13.	Valgrind is programming tool that provides controlled execution, as well as post mortem inspection of an executable.	true/false
14.	The worst-case latency for servicing an interrupt is a combination of factors, including the time taken for higher priority interrupts.	true/false

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15.

static int iSeconds, iMinutes;
void interrupt vUpdateTime(void)
{
++iSeconds;
if (iSeconds>=60) {
iSeconds=0;
++iMinutes;
}
}
) long lSeconds(void)
(
disable();
return (iMinutes*60+iSeconds);
enable();
}
-

Despite disabling interrupts the above pseudo code fails to solve the shared-data problem. true/false

- **16.** An interrupt vector table contains the addresses of the interrupt service routines. true/false
- 17. An interrupt can **not** be serviced faster than the time needed to save the context of code running on the processor. true/false
- 18. Critical sections can be guarded by disabling and enabling interrupts.
 interrupts arriving during such a critical section are buffered and handled upon exit.
 true/false
- **19.** Given is the following RTOS (pseudo) code with priority T1 > T2.

```
void T1(void) {
   while (1) {
      OS_Pend(sem1); // event #1 may unblock any time
      f(1);
      OSTimeDly(1);
   }
}
void T2(void) {
   while (1) \{
      OS_Pend(sem2); // event #2 may unblock any time
      f(-1);
      OSTimeDly(3);
   }
}
void f(int i) {
   OS_Pend(mutex);
   counter = counter + i ; // modify some global counter
   OS_Post (mutex);
}
```

This code suffers from a data sharing problem.

true/false

true/false

20. If the order of events is 1, 2, 1, 2, 1 and they occur within 10 ms from each other, then the final value of the counter will be increased by 1. true/false

21. The function f() is reentrant

22.	The shared-data problem can be solved through enabling interrupts.	true/false
23.	A deadly embrace requires a minimum of 3 tasks of different priority and 1 semaphore to occur.	true/false
24.	When a processor is powered up, interrupts are disabled until further notice.	true/false
25.	While interrupts are disabled atomicity is guaranteed even when calling a non-reentrant function.	true/false
26.	Shared variables marked volatile guarantee atomic access.	true/false
27.	Using interrupts improves system response time.	true/false
28.	The primary shortcoming of an RRI architecture is that it is more complex than RR.	true/false
29.	An RTOS architecture supports priority-based ISRs.	true/false
30.	With an FQS architecture, the worst response time of a task includes the time taken by the longest task in the system.	true/false
31.	With an RTOS every task needs its own stack.	true/false
32.	An RR architecture is most robust to code changes.	true/false
33.	With an RTOS it is impossible to make direct use of harware timers.	true/false
34.	In an RTOS, tasks can be in state BLOCKED, READY or RUNNING. - a task starts in the state RUNNING.	true/false
35.	An ISR may change a task's status from BLOCKED to READY.	true/false
36.	A high-priority task must not invoke an RTOS function that may block.	true/false
37.	When using an RTOS signaling between ISRs and tasks must be done by calling appropriate RTOS primitives.	true/false
38.	A program running on an RTOS may create tasks dynamically at runtime. - the number of tasks is limited by the number of priority levels supported.	true/false
39.	An RTOS usually provides two types of delay functions: polling-based and timer-based. - timer-based delays are more efficient as other tasks can run while the caller is waiting for the specified time to pass.	true/false
40.	Assume that one system clock tick = 10 ms. - Calling the function OSTimeDly(6) causes a delay between 50 and 70 ms.	true/false
41.	To address the shared-data problem, many RTOSs provide communication primitives like queues, mailboxes, and pipes. - a common advantage is that they allow pointers to be passed from one task to another.	true/false
42.	A disadvantage of queues over pipes is that messages/items are handled strictly in FIFO order.	
		true/false

- 43. With the X32 RTOS creating a task amounts to initializing a stack and invoking a context switch to the task's main function.
 This approach provides the possibility to use one stack for multiple (concurrent) tasks and reduce the memory footprint.
- **44.** The **heartbeat timer** is a single hardware timer an RTOS is using to monitor the liveness of the task set involved.

```
45.
    Consider the following code fragment:
    #include <stdio.h>
 1
 2
    #include <string.h>
 3
    #include <stdlib.h>
 4
 5
                                      // copied from <uart.h> for reference
    extern char *UART_rx_buf;
 6
    extern char *UART_tx_buf;
 7
    extern char *UART_ier;
 8
 9
    #define LEN 80
 10
   static char *next_command = NULL;
 11
 12 void rx_ready() {
 13
         static char buffer[2][LEN];
 14
         static int toggle = 0;
 15
         static char *command = buffer[0];
 16
         static int cnt = 0;
 17
         char c = *UART_rx_buf;
 18
 19
         if (c == '\n') {
20
             command[cnt] = ' \setminus 0';
21
             next_command = command;
 22
            toggle = 1 - toggle;
             command = buffer[toggle];
 23
 24
             cnt = 0;
 25
         } else {
26
             command[cnt++] = c;
27
         }
 28
    }
 29
 30
    int main() {
31
         *UART_ier |= 0x3;
                                       // start RX and TX please
32
         while (1) {
33
             if (next_command != NULL) {
 34
                 if (strcmp(next_command, "exit") == 0) {
 35
                      exit(0);
 36
                 } else if (strcmp(next_command, "hello") == 0) {
37
                     printf("world\n");
38
                 }
39
                 next_command = NULL;
40
             }
41
             . . .
 42
         }
 43
    }
```

This code is an example of an RRI architecture.

true/false

true/false

46.	Consider lines 5-7 in which some of a UART's registers are declared. This way a UART, or any other peripheral for that matter, can be accessed with normal read/write instructions.	
	- this mode of operation is called 'memory-mapped I/O'.	true/false
47.	The function rx_ready() uses a technique called 'alternating buffers'. - From line 13 we can infer that the buffers are allocated on the call stack.	true/false
48.	The code suffers from a (subtle) data sharing bug as both rx_ready() and main() write to the same global variable next_command. - in certain cases rx_ready() will overwrite buffered data still to be read by main().	true/false
49.	Removing the write statement on line 39 will not resolve the shared data bug. - it will cause main () to repeat the same command until rx_ready() is invoked again.	true/false
50.	An alternative approach would be to make use of semaphores to support rx_ready() passing the next command to main().	
	- only a single semaphore initialized to 0 is needed.	true/false
51.	Time slicing between tasks of equal priority is to be avoided as it compromises the predictability of their response times.	true/false
52.	The minimal memory footprint of a program grows linearly with the number of tasks.	true/false
53.	A semaphore S used by tasks A and B must be initialized by either A or B.	true/false
54.	An advantage of using tasks is that it allows for better data encapsulation.	true/false
55.	Tasks should have different priorities to avoid fairness issues imposed by the RTOS.	true/false
56.	When developing code for an embedded system, the software can de structured into HW-	
	dependent and HW-independent code. - Doing so makes debugging HW-independent code feasible on the target platform	true/false
57.	A logic analyzer is preferred to an in-circuit emulator because it can be used with any type of processor.	true/false
58.	Although the assert macro is a useful debugging aid during program development, it can only be used on the host.	true/false
59.	A large study of outdoor sensor-network deployments [Beutel:2009] has shown that the water-proof packaging of the base station is key to establishing a reliable connection to the back bone.	true/false
60.	 When debugging code for a distributed sensor network, collecting the (debug) output of the nodes can be arranged in different ways. - A wireless testbed requires no physical instrumentation (i.e. wiring) of the sensor nodes. 	true/false